CS 188: Artificial Intelligence
Spring 2010

Lecture 2: Queue-Based Search
1/21/2010

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Many slides from Dan Klein

Announcements

- Project 0: Python Tutorial
  - Out today. Due next week Thursday.
  - Lab sessions in 271 Soda:
    - Monday 2-3pm
    - Wednesday 4-5pm
  - The lab time is optional, but P0 itself is not
  - On submit, you should get email from the autograder
  - Potentially more lab sessions or office hours held in the lab --- track the announcements section on the webpage!
- Written 1: Search
  - Out today, also due next week Thursday.
- Sections starting next week, location: 285 Cory
  - Section 101: Tue 3-4pm
  - Section 104: Tue 4-5pm
  - Section 102: Wed 11-noon
  - Section 103: Wed noon-1pm

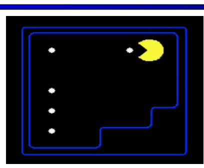
Today

- Agents
- Search Problems
- Uninformed Search Methods (part review for some)
  - Depth-First Search
  - Breadth-First Search
  - Uniform-Cost Search
- Heuristic Search Methods (new for all)
  - Greedy Search

Reminder

- Only a very small fraction of AI is about making computers play games intelligently
- Recall: computer vision, natural language, robotics, machine learning, computational biology, etc.
- That being said: games tend to provide relatively simple example settings which are great to illustrate concepts and learn about algorithms which underlie many areas of AI

A reflex agent for pacman

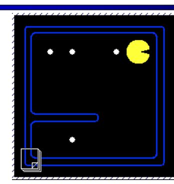


4 actions: move North, East, South or West

Reflex agent

- While(food left)
  - Sort the possible directions to move according to the amount of food in each direction
  - Go in the direction with the largest amount of food

A reflex agent for pacman (2)



Reflex agent

- While(food left)
  - Sort the possible directions to move according to the amount of food in each direction
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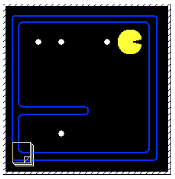
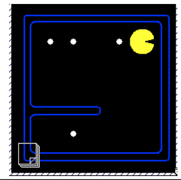
A reflex agent for pacman (3)



Reflex agent

- While(food left)
  - Sort the possible directions to move according to the amount of food in each direction
  - Go in the direction with the largest amount of food
    - But, if other options are available, exclude the direction we just came from

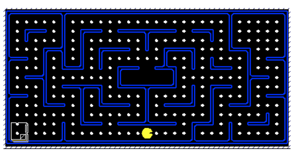
A reflex agent for pacman (4)



Reflex agent

- While(food left)
  - If can keep going in the current direction, do so
  - Otherwise:
    - Sort directions according to the amount of food
    - Go in the direction with the largest amount of food
    - But, if other options are available, exclude the direction we just came from

A reflex agent for pacman (5)



Reflex agent

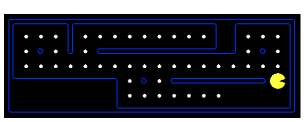
- While(food left)
  - If can keep going in the current direction, do so
  - Otherwise:
    - Sort directions according to the amount of food
    - Go in the direction with the largest amount of food
    - But, if other options are available, exclude the direction we just came from

Reflex Agents

- Reflex agents:
  - Choose action based on current percept (and maybe memory)
  - May have memory or a model of the world's current state
  - Do not consider the future consequences of their actions
  - Act on how the world IS
- Can a reflex agent be rational?

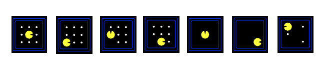
Goal Based Agents

- Goal-based agents:
  - Plan ahead
  - Ask "what if"
  - Decisions based on (hypothesized) consequences of actions
  - Must have a model of how the world evolves in response to actions
  - Act on how the world WOULD BE

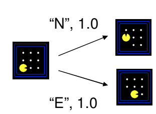


Search Problems

- A search problem consists of:
  - A state space

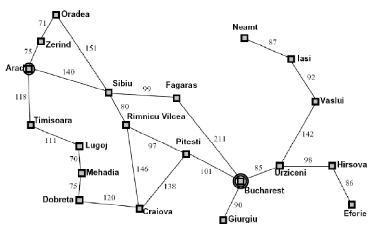


  - A successor function



  - A start state and a goal test
- A solution is a sequence of actions (a plan) which transforms the start state to a goal state

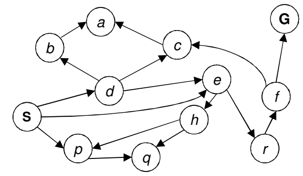
Example: Romania



- State space:
  - Cities
- Successor function:
  - Go to adj city with cost = dist
- Start state:
  - Arad
- Goal test:
  - Is state == Bucharest?
- Solution?

State Space Graphs

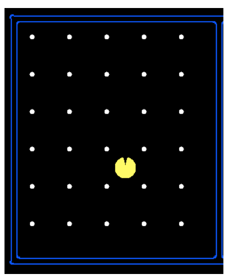
- State space graph: A mathematical representation of a search problem
  - For every search problem, there's a corresponding state space graph
  - The successor function is represented by arcs
- We can rarely build this graph in memory (so we don't)



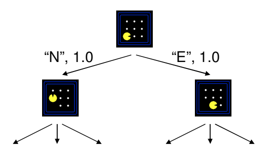
Ridiculously tiny search graph for a tiny search problem

State Space Sizes?

- Search Problem:
  Eat all of the food
- Pacman positions:
  10 x 12 = 120
- Food count: 30

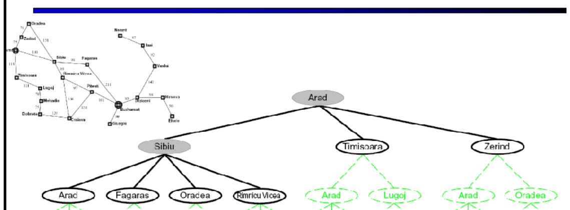


Search Trees



- A search tree:
  - This is a "what if" tree of plans and outcomes
  - Start state at the root node
  - Children correspond to successors
  - Nodes contain states, correspond to PLANS to those states
  - For most problems, we can never actually build the whole tree

Another Search Tree



- Search:
  - Expand out possible plans
  - Maintain a fringe of unexpanded plans
  - Try to expand as few tree nodes as possible

General Tree Search

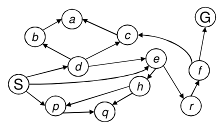
```
function TREE-SEARCH( problem, strategy) returns a solution, or failure
   initialize the search tree using the initial state of problem
   loop do
      if there are no candidates for expansion then return failure
      choose a leaf node for expansion according to strategy
      if the node contains a goal state then return the corresponding solution
      else expand the node and add the resulting nodes to the search tree
   end
```

- Important ideas:
  - Fringe
  - Expansion
  - Exploration strategy

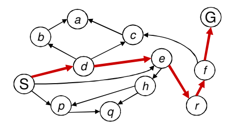
Detailed pseudocode is in the book!

- Main question: which fringe nodes to explore?

Example: Tree Search

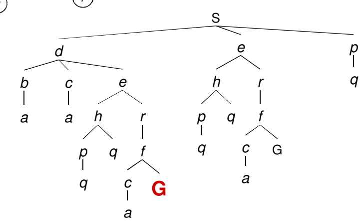


State Graphs vs. Search Trees



Each NODE in in the search tree is an entire PATH in the problem graph.

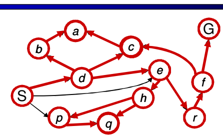
We construct both on demand – and we construct as little as possible.

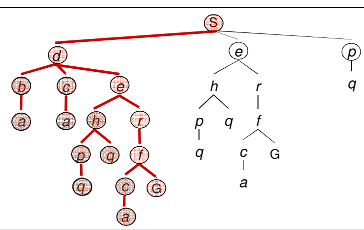


Review: Depth First Search

Strategy: expand deepest node first

Implementation: Fringe is a LIFO stack

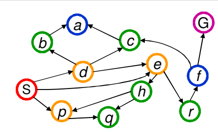




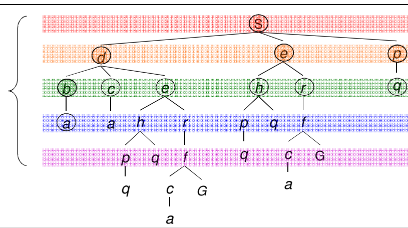
Review: Breadth First Search

Strategy: expand shallowest node first

Implementation: Fringe is a FIFO queue



Search Tiers



Search Algorithm Properties

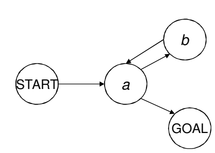
- Complete? Guaranteed to find a solution if one exists?
- Optimal? Guaranteed to find the least cost path?
- Time complexity?
- Space complexity?

Variables:

n	Number of states in the problem
b	The average branching factor B (the average number of successors)
$C*$	Cost of least cost solution
s	Depth of the shallowest solution
m	Max depth of the search tree

DFS

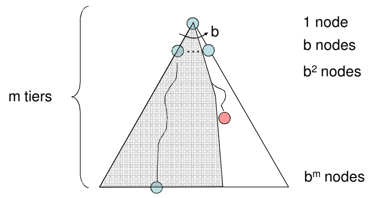
Algorithm		Complete	Optimal	Time	Space
DFS	Depth First Search	N	N	Infinite	Infinite



- Infinite paths make DFS incomplete…
- How can we fix this?

DFS

- With cycle checking, DFS is complete.*



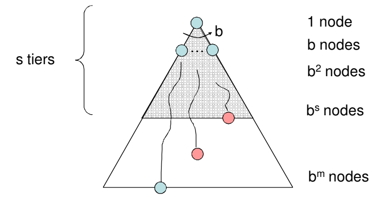
Algorithm		Complete	Optimal	Time	Space
DFS	w/ Path Checking	Y	N	$O(b^m)$	$O(bm)$

- When is DFS optimal?

* Or graph search – next lecture.

BFS

Algorithm		Complete	Optimal	Time	Space
DFS	w/ Path Checking	Y	N	$O(b^{m+1})$	$O(bm)$
BFS		Y	N*	$O(b^{s+1})$	$O(b^{s+1})$



- When is BFS optimal?

Comparisons

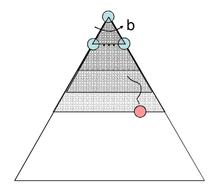
- When will BFS outperform DFS?
- When will DFS outperform BFS?

Iterative Deepening

Iterative deepening uses DFS as a subroutine:

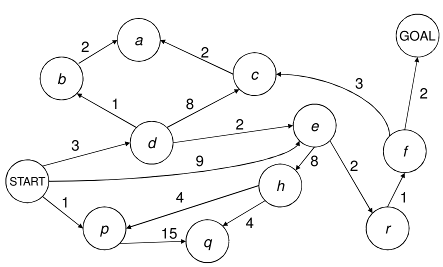
1. Do a DFS which only searches for paths of length 1 or less.
2. If "1" failed, do a DFS which only searches paths of length 2 or less.
3. If "2" failed, do a DFS which only searches paths of length 3 or less.

….and so on.



Algorithm		Complete	Optimal	Time	Space
DFS	w/ Path Checking	Y	N	$O(b^m)$	$O(bm)$
BFS		Y	N*	$O(b^{s+1})$	$O(b^{s+1})$
ID		Y	N*	$O(b^{s+1})$	$O(bs)$

Costs on Actions

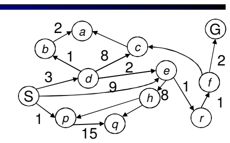


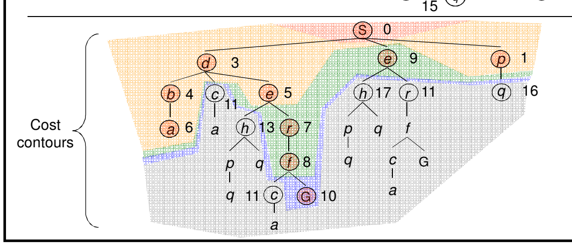
Notice that BFS finds the shortest path in terms of number of transitions. It does not find the least-cost path.
We will quickly cover an algorithm which does find the least-cost path.

Uniform Cost Search

Expand cheapest node first:

Fringe is a priority queue





Priority Queue Refresher

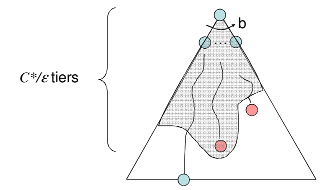
- A priority queue is a data structure in which you can insert and retrieve (key, value) pairs with the following operations:

pq.push(key, value)	inserts *(key, value)* into the queue.
pq.pop()	returns the key with the lowest value, and removes it from the queue.

- You can decrease a key's priority by pushing it again
- Unlike a regular queue, insertions aren't constant time, usually O(log *n*)
- We'll need priority queues for cost-sensitive search methods

Uniform Cost Search

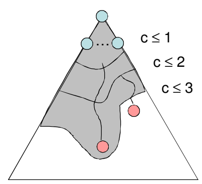
Algorithm		Complete	Optimal	Time	Space
DFS	w/ Path Checking	Y	N	$O(b^m)$	$O(bm)$
BFS		Y	N	$O(b^{s+1})$	$O(b^{s+1})$
UCS		Y*	Y	$O(b^{C*/\varepsilon})$	$O(b^{C*/\varepsilon})$

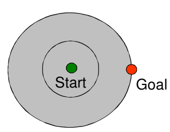


** UCS can fail if actions can get arbitrarily cheap*

Uniform Cost Issues

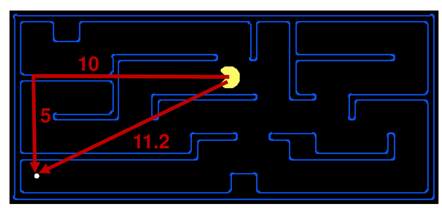
- Remember: explores increasing cost contours
- The good: UCS is complete and optimal!
- The bad:
  - Explores options in every "direction"
  - No information about goal location



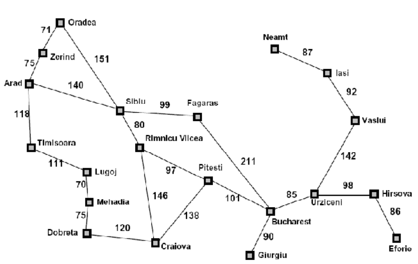


Search Heuristics

- Any *estimate* of how close a state is to a goal
- Designed for a particular search problem
- Examples: Manhattan distance, Euclidean distance



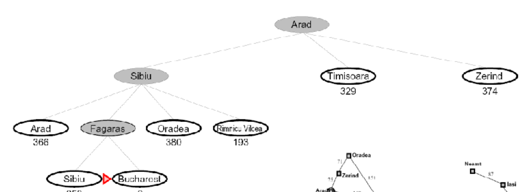
Heuristics



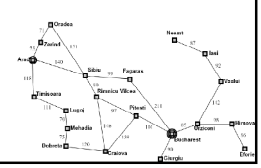
Straight-line distance to Bucharest	
Arad	366
Bucharest	0
Craiova	160
Dobreta	242
Eforie	161
Fagaras	178
Giurgiu	77
Hirsova	151
Iasi	226
Lugoj	244
Mehadia	241
Neamt	234
Oradea	380
Pitesti	98
Rimnicu Vilcea	193
Sibiu	253
Timisoara	329
Urziceni	80
Vaslui	199
Zerind	374

Best First / Greedy Search

- Expand the node that seems closest…

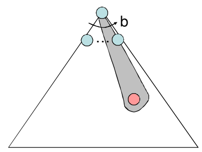


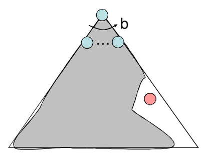
- What can go wrong?



Best First / Greedy Search

- A common case:
  - Best-first takes you straight to the (wrong) goal
- Worst-case: like a badly-guided DFS in the worst case
  - Can explore everything
  - Can get stuck in loops if no cycle checking
- Like DFS in completeness (finite states w/ cycle checking)





- Can we leverage the heuristic information in a more sound way?

→A* search

We will cover that on Tuesday!
